

# Some definitions for the Cut-ELM. Thm.

• degree of a formula  $A$ , (notation:  $|A|$ ), is the number of logical connectives in  $A$ .

• cut-rank of a proof  $\pi$  in  $L_k$ , (notation:  $c(\pi)$ ) :

-  $\pi$  axiom, then  $c(\pi) = 0$

-  $\pi$  has shape 
$$p \frac{\frac{\nabla}{\Sigma'}}{\Sigma} : c(\pi) = c(\pi')$$

-  $\pi$  has shape 
$$p \frac{\frac{\nabla}{\Sigma'} \quad \frac{\nabla}{\Sigma''}}{\Sigma} :$$

• if  $p \neq \text{cut}$  then  $c(\pi) = \max\{c(\pi'_*), c(\pi'')\}$

• if  $p = \text{cut}$  and  $A$  is the cut-formula then  $c(\pi) = \max\{c(\pi'), c(\pi''), |A| + 1\}$

• depth of a proof tree  $\pi$  (notation  $d(\pi)$ ) :

-  $\pi$  axiom, then  $d(\pi) = 0$

-  $\pi$  has shape 
$$p \frac{\frac{\nabla}{\Sigma'}}{\Sigma} \text{ then } d(\pi) = d(\pi') + 1$$

-  $\pi$  has shape 
$$p \frac{\frac{\nabla}{\Sigma'} \quad \frac{\nabla}{\Sigma''}}{\Sigma} \text{ then } d(\pi) = \max\{d(\pi'_*), d(\pi''),$$

$d(\pi'')\} + 1.$

(• logical depth: similar to depth, but structural rules are not counted).

# REDUCTION LEMMA (Tait-Girard)

Given proofs  $\frac{\pi_1}{\Gamma \vdash \Delta, mA}$   $\frac{\pi_2}{uA, \Lambda \vdash \Theta}$  in  $Lk^+$

where  $m, u > 0$  (assume  $c(\pi_1), c(\pi_2) \leq |A|$ )

We can construct a proof in  $Lk^+$

$\frac{\pi}{\Gamma, \Lambda \vdash \Delta, \Theta}$  such that  $c(\pi) \leq |A|$ .

(Moreover  $e(\pi) \leq 2(e(\pi_1) + e(\pi_2))$  and if the rules of  $\supset$  are omitted then  $e(\pi) = e(\pi_1) + e(\pi_2)$ .)

proof by induction on  $d(\pi_1) + d(\pi_2)$  when

$\pi_1, \pi_2$  are the immediate subtrees of the proof tree

or  $\frac{\frac{\pi_1}{\Gamma \vdash \Delta, (mA)} \quad \frac{\pi_2}{(uA), \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$  A cut formula.

Many cases, not mutually exclusive one.

Organisation must be careful!

The transformed proof is  $\pi$ .

1<sup>st</sup> Level of organisation : 4 cases

2<sup>nd</sup> level : L/R + simple/multiple cut.

① The root  $\Pi_1$  and the root  $\Pi_2$  are conclusions of some LOGICAL inference, having some occurrences of  $A$ , the cut formula, as principal formula. In other words  $A$  is ACTIVE in the logical inference.

We must consider two versions:

- simple cut
- cross cuts

② Either the root of  $\Pi_1$  or of  $\Pi_2$  is the conclusion of some logical rule, or of a cut, or of a structural rule, having  $X \neq A$  as principal formula:  $A$  is PASSIVE

We will consider the case it happens in the left subtree; on the right it is symmetric.

③ Either  $\Pi_1$  or  $\Pi_2$  is an axiom.

We will consider the left subtree; the right is symmetric

④ Either the root of  $\Pi_1$  or of  $\Pi_2$  is the conclusion of some weakening or contraction on the cut formula  $A$  -

- We will consider the left subtree (right is symm)
- Distinguish simple/crosscut.

① A ACTIVE in some logical rule TC3  
(1.1.)

(i)  $\wedge_R \parallel \wedge_L$

$$\begin{array}{c}
 \begin{array}{c}
 \pi_1 \quad \pi_2 \\
 \Gamma \vdash \Delta, B \quad \Gamma \vdash \Delta, C \\
 \hline
 \Gamma \vdash \Delta, B \wedge C \\
 \wedge_R \\
 \hline
 \Gamma, \lambda \vdash \Delta, \Theta
 \end{array}
 \quad
 \begin{array}{c}
 \pi_3 \\
 B, \lambda \vdash \Theta \\
 \hline
 B \wedge C, \lambda \vdash \Theta \\
 \wedge_L \\
 \hline
 \Gamma, \lambda \vdash \Delta, \Theta
 \end{array}
 \quad
 \Rightarrow
 \quad
 \begin{array}{c}
 \pi_1 \quad \pi_3 \\
 \Gamma \vdash \Delta, B \quad B, \lambda \vdash \Theta \\
 \hline
 \Gamma, \lambda \vdash \Delta, \Theta \\
 \wedge
 \end{array}
 \end{array}$$

hp:  $c(\pi_1), c(\pi_2) \leq |A|$

where  $A = B \wedge C$

$|A| \geq |B| + 1$  and  $c(\pi_1), c(\pi_2) \leq c(\pi_1)$   
 $c(\pi_3) \leq c(\pi_2)$

hence  $c(\pi) \leq |A|$

because  $c(\pi) = \max \{ |B| + 1, c(\pi_1), c(\pi_2) \}$

Cross-cuts: assume  $m, n > 0$  (with  $m$  or  $n = 0$  there are simplifications).

let  $A = B \wedge C$

$$\begin{array}{c}
 \begin{array}{c}
 \pi_1 \quad \pi_2 \\
 \Gamma \vdash \Delta, mA, B \quad \Gamma \vdash \Delta, mA, C \\
 \hline
 \Gamma \vdash \Delta, (m+1)A \\
 \wedge_R \\
 \hline
 \Gamma, \lambda \vdash \Delta, \Theta
 \end{array}
 \quad
 \begin{array}{c}
 \pi_3 \\
 B, nA, \lambda \vdash \Theta \\
 \hline
 (n+1)A, \lambda \vdash \Theta \\
 \wedge \\
 \hline
 \Gamma, \lambda \vdash \Delta, \Theta
 \end{array}
 \end{array}$$

Obtain  $\pi_1'$  and  $\pi_2'$  by applying the inductive hp. to the following proofs:

$$\frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA, B} \quad \frac{\pi_3}{\Lambda \vdash \frac{B, uA, \Lambda \vdash \Theta}{(n+1)A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B}}{\text{i.h.} \rightarrow \Pi_1'}$$

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA, B} \quad \frac{\pi_2}{\Gamma \vdash \Delta, mA, C}}{\Lambda \vdash \frac{\Gamma \vdash \Delta, (m+1)A}{\Gamma \vdash \Delta, (m+1)A}} \quad \frac{\pi_3}{B, uA, \Lambda \vdash \Theta}}{\text{i.h.} \leftarrow \Pi_2' \quad \Gamma, \Lambda \vdash \Delta, \Theta}$$

Let  $\pi$  be

$$\frac{\frac{\frac{\pi_1'}{\Gamma, \Lambda \vdash \Delta, \Theta, B} \quad \frac{\pi_2'}{B, \Gamma, \Lambda \vdash \Delta, \Theta}}{\Gamma, \Gamma, \Lambda, \Lambda \vdash \Delta, \Delta, \Theta, \Theta}}{\text{(str. rules)} \quad \Gamma, \Lambda \vdash \Delta, \Theta}$$

They hold:

$$\left. \begin{array}{l} c(\pi_1) \leq c(\Pi_1) \\ c(\pi_2) \leq c(\Pi_1) \\ c(\pi_3) \leq c(\Pi_2) \end{array} \right\} \xrightarrow{\text{i.h.}} \begin{array}{l} c(\Pi_1'), c(\Pi_2') \in |A| \text{ and} \\ \text{since } |A| \geq |B|+1 \text{ (} A=B \wedge C \text{)} \\ c(\pi) = \max\{|B|+1, c(\Pi_1'), \\ c(\Pi_2')\} \end{array}$$



(ii)  $V_R/U_L$

$$\alpha \frac{\frac{\pi_1}{\Gamma \vdash \Delta, B} \quad \frac{\frac{\pi_2}{B, \Lambda \vdash \Theta} \quad \frac{\pi_3}{C, \Lambda \vdash \Theta}}{B \vee C, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta} \longrightarrow \alpha \frac{\frac{\pi_1}{\Gamma \vdash \Delta, B} \quad \frac{\pi_2}{B, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

$c(\pi_1), c(\pi_2) \leq |A|$  and  $c(\pi) \leq |A|$  because  
 $c(\pi) = \max\{|B|+1, c(\pi_1), c(\pi_2)\}$

and  $c(\pi_1) \leq c(\pi_1)$   
 $c(\pi_2) \leq c(\pi_2)$   
 $c(\pi_3) \leq c(\pi_2)$

crosscut

similar to the previous case (i). Try.

(iii)  $\supset_L/\supset_R$ : Do it as exercise

(iv)  $\neg_R/\neg_R$ :

$$\alpha \frac{\frac{\frac{\pi_1}{A, \Gamma \vdash \Delta} \quad \frac{\pi_2}{\Lambda \vdash \Theta, A}}{\neg A, \Lambda \vdash \Theta} \quad \frac{\pi_1}{\Gamma \vdash \Delta, \neg A}}{\Gamma, \Lambda \vdash \Delta, \Theta} \longrightarrow \alpha \frac{\frac{\pi_2}{\Lambda \vdash \Theta, A} \quad \frac{\pi_1}{\neg A, \Gamma \vdash \Delta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

$c(\pi_1), c(\pi_2) \leq |\neg A| \longrightarrow c(\pi) \leq |\neg A|$  since  
 $c(\pi) = \max\{|\neg A|+1, c(\pi_1), c(\pi_2)\}$  and  
 $c(\pi_1) \leq c(\pi_1)$        $c(\pi_2) \leq c(\pi_2)$ .

concluts:

1.4  
1.4

$$\begin{array}{c} \pi_1 \\ \hline A, \Gamma \vdash \Delta, m(\neg A) \\ \hline \Gamma \vdash \Delta, (m+1)(\neg A) \\ \infty \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta \end{array} \quad \begin{array}{c} \pi_2 \\ \hline m(\neg A), \Lambda \vdash \Theta, A \\ \hline (m+1)(\neg A), \Lambda \vdash \Theta \end{array}$$

Take  $\pi_1'$ ,  $\pi_2'$  obtained by applying the i.h. to, resp.:

$$\begin{array}{c} \pi_1 \\ \hline A, \Gamma \vdash \Delta, m(\neg A) \\ \hline \Gamma \vdash \Delta, (m+1)(\neg A) \\ \infty \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, A \end{array} \quad \begin{array}{c} \pi_2 \\ \hline m(\neg A), \Lambda \vdash \Theta, A \\ \hline (m+1)(\neg A), \Lambda \vdash \Theta \end{array} \rightsquigarrow \pi_1'$$

$$\pi_2' \longleftarrow \begin{array}{c} \pi_1 \\ \hline A, \Gamma \vdash \Delta, m(\neg A) \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, A \\ \infty \\ \hline A, \Gamma, \Lambda \vdash \Delta, \Theta \end{array} \quad \begin{array}{c} \pi_2 \\ \hline m(\neg A), \Lambda \vdash \Theta, A \\ \hline (m+1)(\neg A), \Lambda \vdash \Theta \end{array}$$

and take  $\pi$  as:

$$\begin{array}{c} \pi_1' \quad \pi_2' \\ \Gamma, \Lambda \vdash \Delta, \Theta, A \quad A, \Gamma, \Lambda \vdash \Delta, \Theta \\ \hline \Gamma, \Gamma, \Lambda, \Lambda \vdash \Delta, \Delta, \Theta, \Theta \\ \hline \hline \Gamma, \Lambda \vdash \Delta, \Theta \end{array}$$

$$c(\pi_1) \leq c(\pi_1')$$

$$c(\pi_2) \leq c(\pi_2')$$

$c(\pi_1')$ ,  $c(\pi_2')$   $\leq |\neg A|$  and  $c(\pi) \leq |\neg A|$  because

$$c(\pi) = \max\{|\neg A|+1, c(\pi_1'), c(\pi_2')\}.$$

(V)  $\forall R / \forall R$

TG7  
1.5

$$\frac{\forall R \frac{\frac{\pi_1}{\Gamma \vdash \Delta, B[y/x]}}{\Gamma \vdash \Delta, \forall x.B} \quad \forall L \frac{\frac{\pi_2}{B[t/x], \Lambda \vdash \Theta}}{\forall x.B, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta} \rightarrow \forall R \frac{\frac{\pi_1[t/y]}{\Gamma \vdash \Delta, B[t/x]} \quad \pi_2}{\Gamma, \Lambda \vdash \Delta, \Theta} \frac{B[t/x], \Lambda \vdash \Theta}{\Gamma \vdash \Delta, \Theta}}$$

it might be necessary to rename some eigenvariables in  $\pi_1$  so that they are distinct from all variables in  $t$ .

$$c(\pi_1), c(\pi_2) \leq |\forall x.B| \quad c(\pi) \leq |\forall x.B| \text{ since}$$

$$c(\pi) = \max \{ |B[t/x]| + 1, c(\pi_1[t/y]), c(\pi_2) \}$$

$$c(\pi_1[t/y]) = c(\pi_1) \leq c(\pi_1)$$

$$c(\pi_2) \leq c(\pi_2)$$

cross cut:

$$\frac{\forall R \frac{\frac{\pi_1}{\Gamma \vdash \Delta, u(\forall x.B), B[y/x]}}{\Gamma \vdash \Delta, (u+1)(\forall x.B)}}{\Gamma, \Lambda \vdash \Delta, \Theta} \quad \forall L \frac{\frac{\pi_2}{B[t/x], u(\forall x.B), \Lambda \vdash \Theta}}{(u+1)(\forall x.B), \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

let  $\pi_1', \pi_2'$  be obtained by applying the i.h. to these:

$$\frac{\forall R \frac{\frac{\pi_1}{\Gamma \vdash \Delta, u(\forall x.B), B[y/x]}}{\Gamma, \Lambda \vdash \Delta, \Theta, B[y/x]} \quad \forall L \frac{\frac{\pi_2}{B[t/x], u(\forall x.B), \Lambda \vdash \Theta}}{(u+1)(\forall x.B), \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B[y/x]} \rightsquigarrow \pi_1'$$

and

$$\frac{\forall R \frac{\frac{\pi_1}{\Gamma \vdash \Delta, u(\forall x.B), B[y/x]}}{\Gamma \vdash \Delta, (u+1)\forall x.B} \quad \forall L \frac{\frac{\pi_2}{B[t/x], \Gamma, \Lambda \vdash \Delta, \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}}{B[t/x], \Gamma, \Lambda \vdash \Delta, \Theta} \rightsquigarrow \pi_2'$$

let  $\pi$  be

$\pi_8$   
1.6

$$\begin{array}{l} \pi_1' [t/y] \qquad \qquad \qquad \pi_2' \\ P, \lambda \vdash \Delta, \Theta, B[t/x] \qquad B[t/x], P, \lambda \vdash \Delta, \Theta \\ \hline P, P, \lambda, \lambda \vdash \Delta, \Delta, \Theta, \Theta \\ \hline P, \lambda \vdash \Delta, \Theta \end{array}$$

it might be necessary to rename some  
existential variables in  $\pi_1'$  to make them  
distinct from all variables in  $\pi_2$ .

$$\left. \begin{array}{l} c(\pi_1) \leq c(\pi_1) \\ c(\pi_2) \leq c(\pi_2) \\ c(\pi_1'), c(\pi_2') \leq |Vx.B| \end{array} \right\} \rightarrow c(\pi) \leq |Vx.B|$$

because  $c(\pi) = \max \{ |B[t/x]|, c(\pi_1' [t/y]), c(\pi_2') \}$   
and  $c(\pi_1' [t/y]) = c(\pi_1')$ .

(vi)  $\exists R / \exists L$

$$\begin{array}{l} \pi_1 \qquad \qquad \qquad \pi_2 \qquad \qquad \text{some} \\ P \vdash \Delta, B[t/x] \qquad B[t/y/x], \lambda \vdash \Theta \qquad \text{renaming} \\ \exists R \frac{}{P \vdash \Delta, \exists x.B} \qquad \exists L \frac{}{\exists x.B, \lambda \vdash \Theta} \\ \hline P, \lambda \vdash \Delta, \Theta \qquad \rightarrow \\ \pi_1 \qquad \qquad \pi_2 [t/y] \\ P \vdash \Delta, B[t/x] \qquad B[t/x], \lambda \vdash \Theta \\ \hline P, \lambda \vdash \Delta, \Theta \end{array}$$

non-cut: exercise

② A is passive (on the left); any rule (right) | T49  
2.1

symmetric case if A is on the right subtree, or left logical / str. rules are applied.

(i)  $V_R$ :

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA, B} \quad \frac{\pi_2}{\mu A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \vee C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \vee C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \vee C} \rightarrow \frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA, B} \quad \frac{\pi_2}{\mu A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \vee C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \vee C}$$

$c(\pi_1) \leq c(\pi_1)$   $c(\pi_2) \leq c(\pi_2)$  apply the i.h. on  $\Gamma, \Lambda \vdash \Delta, \Theta, B$ .

(ii)  $\wedge_R$ :

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, \mu A, B} \quad \frac{\pi_2}{\Gamma \vdash \Delta, \mu A, C}}{\Gamma \vdash \Delta, \mu A, B \wedge C} \quad \frac{\pi_3}{\mu A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \wedge C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \wedge C} \rightarrow$$

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, \mu A, B} \quad \frac{\pi_3}{\mu A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, B} \quad \frac{\frac{\pi_2}{\Gamma \vdash \Delta, \mu A, C} \quad \frac{\pi_3}{\mu A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta, C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \wedge C}}{\Gamma, \Lambda \vdash \Delta, \Theta, B \wedge C}$$

$c(\pi_1), c(\pi_2) \leq c(\pi_1)$

$c(\pi_3) \leq c(\pi_2)$

Apply the i.h. to the subtrees rooted at  $\Gamma, \Lambda \vdash \Delta, \Theta, B$

and  $\Gamma, \Lambda \vdash \Delta, \Theta, C$

(iii)  $\supset_R$ : exercise

(iv)  $\neg_R$ :

$$\begin{array}{ccc} \begin{array}{c} \pi_1 \\ \hline \Gamma, \Gamma \vdash \Delta, \neg A \\ \hline \Gamma \vdash \Delta, \neg A, \neg B \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, \neg B \end{array} & \begin{array}{c} \pi_2 \\ \hline \neg A, \Lambda \vdash \Theta \end{array} & \begin{array}{ccc} \pi_1 & & \pi_2 \\ \hline \Gamma, \Gamma \vdash \Delta, \neg A & & \neg A, \Lambda \vdash \Theta \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, \neg B & \xrightarrow{\quad} & \Gamma, \Lambda \vdash \Delta, \Theta, \neg B \end{array} \end{array}$$

$c(\pi_1) \leq c(\pi_2)$   
 $c(\pi_2) \leq c(\pi_2)$  ; apply the i.h. to  $\Gamma, \Gamma, \Lambda \vdash \Delta, \Theta$

(v)  $\forall_R$ :

$$\begin{array}{ccc} \begin{array}{c} \pi_1 \\ \hline \Gamma \vdash \Delta, \neg A, B[y/x] \\ \hline \Gamma \vdash \Delta, \neg A, \forall x. B \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, \forall x. B \end{array} & \begin{array}{c} \pi_2 \\ \hline \neg A, \Lambda \vdash \Theta \end{array} & \rightarrow \end{array}$$

$$\begin{array}{ccc} \begin{array}{c} \pi_1 [z/x] \\ \hline \Gamma \vdash \Delta, \neg A, B[z/y] \\ \hline \Gamma, \Lambda \vdash \Delta, \Theta, B[z/x] \\ \hline \forall_R \Gamma, \Lambda \vdash \Delta, \Theta, \forall x. B \end{array} & \begin{array}{c} \pi_2 \\ \hline \neg A, \Lambda \vdash \Theta \end{array} & \end{array}$$

$c(\pi_1 [z/y]) = c(\pi_1) \leq c(\pi_2)$   
 $c(\pi_2) \leq c(\pi_2)$   
 Apply the i.h. to the subtree rooted  $\Gamma, \Lambda \vdash \Delta, \Theta, B[z/x]$ .

(vi)  $\exists_R$ : exercise

(vii) cont:

TC 11  
2.3

$$\begin{array}{c}
 \pi_1 \\
 \Gamma_1 \vdash \Delta_1, m_1 A, pB \\
 \hline
 \pi_2 \quad \pi_3 \\
 \Gamma_2, \Lambda_1 \vdash \Theta_1, m_2 A \quad mA, \Lambda \vdash \Theta \\
 \hline
 \Gamma \vdash \Delta, mA \\
 \hline
 \Gamma, \Lambda \vdash \Delta, \Theta \quad \rightarrow
 \end{array}$$

Notation:  $m_1 + m_2 = m$   $\Gamma = \Gamma_1, \Lambda_1$   $\Delta = \Delta_1, \Theta_1$

By hp  $c(\pi_1), c(\pi_2) \leq |A|$  and  
 $c(\pi_1) = \max\{|B|+1, c(\pi_1), c(\pi_2)\}$ .

$\Rightarrow |B| < |A| \quad c(\pi_i) \leq |A| \quad i=1..3$

so, in particular  $B \neq A$ .

We consider the transformation for  $m_1, m_2 > 0$   
 (special cases are if one between  $m_1$  ( $m_2 = 0$ )).

$\pi_1'$  is obtained by applying the transformation to

$$\begin{array}{c}
 \pi_1 \\
 \Gamma_1 \vdash \Delta_1, pB, m_1 A \\
 \hline
 \pi_3 \\
 mA, \Lambda \vdash \Theta \\
 \hline
 \Gamma_1, \Lambda \vdash \Delta_1, \Theta, pB
 \end{array}$$

and  $\pi_2'$  similarly from

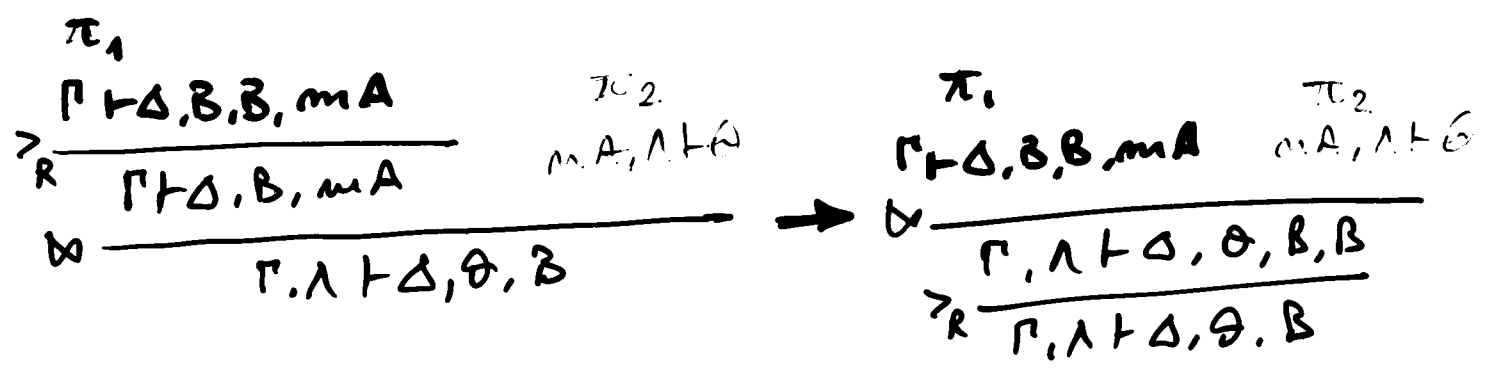
$$\begin{array}{c}
 \pi_2 \\
 \Gamma_2, \Lambda_1 \vdash \Theta_1, m_2 A \\
 \hline
 \pi_3 \\
 mA, \Lambda \vdash \Theta \\
 \hline
 \Gamma_2, \Lambda_1, \Lambda \vdash \Theta_1, \Theta \quad ;
 \end{array}$$

then take

$$\begin{array}{c}
 \pi_1' \\
 \Gamma_1, \Lambda \vdash \Delta_1, \Theta, pB \\
 \hline
 \pi_2' \\
 \Gamma_2, \Lambda_1, \Lambda \vdash \Theta_1, \Theta \\
 \hline
 \Gamma, \Lambda, \Lambda \vdash \Delta, \Theta, \Theta \\
 \hline
 \Gamma, \Lambda \vdash \Delta, \Theta
 \end{array}$$

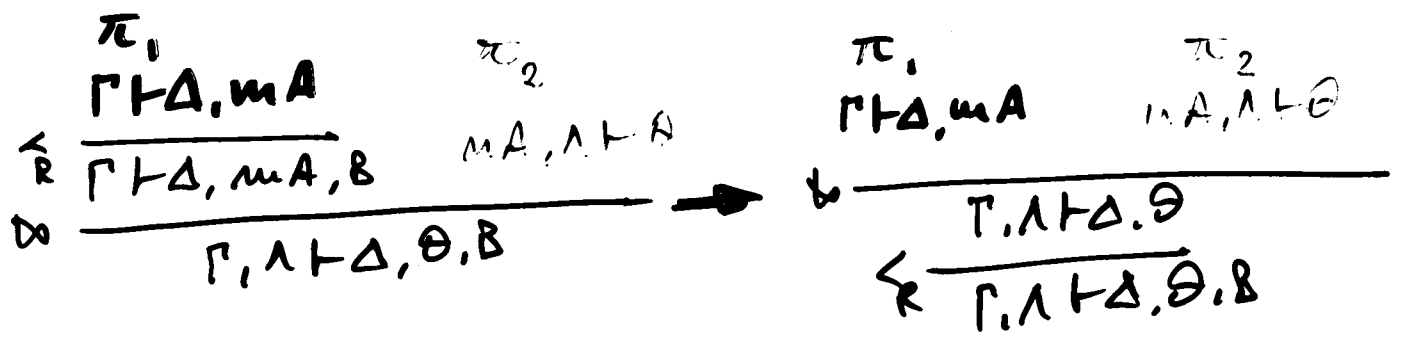
$c(\pi_1') \leq |A|$ ,  
 $c(\pi_2') \leq |A|$ ,  
 $|B| < |A|$   
 $\Rightarrow c(\pi) \leq |A|$ .

(viii)  $\succ_R$



$c(\pi_1) \leq c(\Pi_1)$   
 $c(\pi_2) \leq c(\Pi_2)$ . Apply the i.h. to the subterm noted  $\Gamma, \Lambda \vdash \Delta, \Theta, B$ .

(ix)  $\prec_R$



$c(\pi_1) \leq c(\Pi_1)$   
 $c(\pi_2) \leq c(\Pi_2)$ ; Apply the i.h. to the subterm noted  $\Gamma, \Lambda \vdash \Delta, \Theta$ .

✓

③ Either  $\pi_1$  or  $\pi_2$  is an axiom  
 (we consider the left subtree; the other is symmetric) TG 13  
3.1

$$\begin{array}{c} \text{A} \vdash \text{A} \quad \pi_2 \\ \hline \text{A}, \Lambda \vdash \Theta \end{array} \xrightarrow{\text{L}} \begin{array}{c} \pi_2 \\ \text{A}, \Lambda \vdash \Theta \\ \hline \hline \text{A}, \Lambda \vdash \Theta \end{array}$$

$c(\pi_2) \leq c(\pi_2)$ , by hp  $c(\pi_1), c(\pi_2) \leq |A|$   
 hence  $c(\pi) \leq |A|$ .

④ The root of  $\pi_1$  ( $\pi_2$ ) is the conclusion of some weakening/contraction on the cut formula  $A$ .  
(We consider this to happen on the left subtree,  $\pi_1$ ).

(i)  $<_R$ :

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta} \quad \frac{\pi_2}{A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, A} \quad \frac{\pi_2}{A, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta} \longrightarrow \frac{\frac{\pi_1}{\Gamma \vdash \Delta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

general case,  $m, n > 0$ :

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma \vdash \Delta, (m+1)A} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta} \longrightarrow \frac{\frac{\pi_1}{\Gamma \vdash \Delta, mA} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

By hp:  $c(\pi_1), c(\pi_2) \in |A| \rightarrow c(\pi) \in |A|$  in the first case.  
Otherwise  $c(\pi_1) \in c(\pi_1), c(\pi_2) \in c(\pi_2)$  and use the i.h.

(ii)  $>_R$ :

$$\frac{\frac{\frac{\pi_1}{\Gamma \vdash \Delta, (m-1)A, A, \Delta} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma \vdash \Delta, mA} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta} \longrightarrow \frac{\frac{\pi_1}{\Gamma \vdash \Delta, (m+1)A} \quad \frac{\pi_2}{mA, \Lambda \vdash \Theta}}{\Gamma, \Lambda \vdash \Delta, \Theta}$$

since by hp  $c(\pi_1), c(\pi_2) \in |A|$  and  $c(\pi_i) \in c(\pi_i) \text{ } i=1,2$ ;  
Q.E.D. by the i.h.

Then TAIT '68 provides an upper bound on the size of the resulting cut-free proof.

Given a proof  $\frac{\Delta}{\Gamma \vdash \Delta}$  with cut-rank  $c(\pi)$

there is a cut-free proof  $\frac{\Delta}{\Gamma \vdash \Delta}$  such that

$$l(\pi^*) \leq \exp(4, c(\pi), l(\pi))$$

Then GENTZEN '35 provide a transformation (algorithm) for LK and LJ (inhabitantic logic in the sequent calculus).